## Trees



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Manager-Employee Relation


Google Maps



Planetary Hierarchy


Modeling Computation

## Nomenclature

- Root
- Stem

Edges

- Branches
- Leaves
- Fruits
- Flowers


## Definition

A tree is a collection of nodes.
It could be empty.
Otherwise, it contains a root node, connected to zero or more (child) nodes, each of which is a tree in itself!


Alternatively, a tree is a collection of nodes and directed edges, such that each node except one has a single parent. The node without a parent node is the root.

## Nomenclature

Empty Tree

Root has no parent. Leaves have no children.
Non-leaves are internal nodes.

Tree with one node

Tree with two nodes

Trees with three nodes


## Properties

- A tree has six nodes.
- What is the minimum number of edges in the tree?
- What is the maximum?
- Generalization for N nodes?
- How many (undirected) paths exist between two nodes?


## More Nomenclature

- Sibling
- What is the maximum number of siblings a node may have in an N node tree?
- Grandparent, grandchild
- Ancestor, descendant
- Path, length
- Height, depth



## Exercises

- Given (a pointer to) a node in an employee tree, list all its direct and indirect subordinates.
- Same as above with the name of the employee given.
- Find distance between two nodes.
- Find tree diameter (max. distance).
- Convert infix to postfix (using a tree).
- Mirror a tree vertically.
- Find if there is a directed path from p to q .


## Learning Outcomes

- Apply tree data structure in relevant applications.
- Construct trees in C++ and perform operations such as insert.
- Perform traversals on trees.
- Analyze complexity of various operations.


## Implementation

- A challenge is that the maximum number of children is unknown, and may vary dynamically.

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## Directory Listing

## There is a Linux command to list a directory in a tree-like format. Any guesses for the command?





- 1.cpp
- 2.cpp
- trees.pdf
searchtree
- 1.cpp
bst.pdf
- saurabh

L cv.pdf
somesh


## Switch to code.

2.cpp and 3.cpp

## Traversals

- Preorder
- Process each node before processing its children.
- Children can be processed in any order.
- Postorder
- Process each node after processing its children.
- Children can be processed in any order.
- Preorder and postorder are examples of Depth-First Traversal.
- Children of a node are processed before processing its siblings.
- The other way is called Breadth-First or Level-Order ${ }_{3}$ Traversal.


## Preorder

## Iterative

```
void Tree::preorder() {
    std::stack<PtrToNode> stack;
    stack.push(root);
    while (!stack.empty()) {
        PtrToNode rr = stack.top();
        stack.pop();
        if (rr) {
            rr->print();
            for (auto child: rr->children)
        stack.push(child);
    }
    }
}
```


## Switch to code: 4.cpp, 6.cpp

Classwork: Indent files as per their depth. What is the code complexity? Note that indentation time also needs to be considered.

## Find full size of each directory



## Postorder

## Iterative

Recursive
void Tree:: postorder(PtrToNode rr) \{
if (rr) \{
for (auto child:rr->children)
postorder(child);
rr->print();
\}
}
}
void Tree::postorder() \{
postorder(root);
\}

Try it out offline.

## Switch to code: 5.cpp

## Story so far...

- General trees
- arbitrary number of children
- Resembles several situations such as employees, files, ...
- Special trees
- Fixed / bounded number of children
- Resembles situations such as expressions, boolean flows, ...
- All the children may not be present.


## K-ary Trees

```
typedef struct TreeNode *PtrToNode;
struct TreeNode {
    int data;
    PtrToNode firstChild;
    PtrToNode nextSibling;
};
```


## For a fixed K

typedef struct TreeNode *PtrToNode;
struct TreeNode \{ int data;
PtrToNode children[K];
\};

```
#include <vector>
typedef struct TreeNode *PtrToNode;
struct TreeNode {
    int data;
    std::vector<PtrToNode> children;
};
```


## When K == 2

typedef struct TreeNode *PtrToNode;
struct TreeNode \{ int data;
PtrToNode left;
PtrToNode right;

## K-ary Trees



For a fixed K
typedef struct TreeNode *PtrToNode;
struct TreeNode \{ int data;
PtrToNode children[K];
\};

When K == 2
typedef struct TreeNode *PtrToNode;
struct TreeNode \{ int data;
PtrToNode left; PtrToNode right;
\};

## Properties of Binary Trees

- For an N node binary tree ( $\mathrm{N}>0$ ):
- What is the maximum height? $\mathbf{N}$
- What is the minimum height? $\quad \log _{2}(\mathbf{N})$
- How many NULL pointers? $\quad \mathrm{N}+1$
- How many min/max leaves? 0/1, N/2
- What is the maximum number of nodes a binary tree of height H may have?
- Full nodes (nodes with two children):
- how many minimum, maximum? 0,N/2-1
- Show that \#full nodes $+1==$ \#leaves in a non $\bar{z}_{\overline{2}}$ empty binary tree.


## Operations on Trees

- Insert: our addChild would take care of this.
- Given pointers, this is constant time operation.
- Remove: Update parent's pointer to NULL (and free memory).
- What if the node getting removed has children?
- Based on the above answer, the complexity could be $\mathrm{O}(1)$ or $\mathrm{O}(\mathrm{N})$
- Search: Our tree traversals can help.
- Can a tree contain duplicate values?
- This is $\mathrm{O}(\mathrm{N})$, since the whole tree needs to be searched in the worst case.


## Some Questions?

-What if a child node is common to two parents?

- Ancestry
- Can the edge be undirected?
- Can the edges have weights?
- Can there be multiple roots?
- Can there be multiple edges between two nodes?
- Is it okay to draw a tree with root at the bottom?


## Coding (a little different)

- I want to transmit some data.
- Data contains a-z and space.
- For these 27 characters, I need 5 bits.
- For N characters, I need $\log _{2}(\mathrm{~N})$ bits.
- Encoding pattern:
- space $=00000, a=00001, b=00010, \ldots, z=11010$
- Decoding:
- Each 5-bit string represents a unique character (except the last five strings: 11011 to 11111).
- What is $001001100101110000010110101111 ?$


## How is a code related to a tree?



- It is a binary tree.
- Tree is (almost) complete.
- Has height of 5, equal to the code length.
- Each character has a unique code, because each tree node has a unique path from the root.
- Encoding: Given a character, traverse back from its node towards the root, and we get the reverse of its code.
- Decoding: Given a code, traverse the tree from the root, and the node we reach is the corresponding character.
- None of the interior nodes represents a character.


## What is this plot?



## Make the common case faster!

- If most people order vanilla ice-cream, keep it in front.
- If only a few students buy fish, keep it at a separate counter.
- If most people coming to the department use bicycle, bicycle parking should be prioritized.
- If most of the humans in the classroom stay in hostels, the classes should be held in hostels!
- If 'e' gets used more often, can we transmit it faster?


## Shorter Codes

| Character | Frequency | Code | Code 2 | Code 3 |
| :---: | :---: | :---: | :---: | :---: |
| e | 12.02 | 0 | 0 | 00 |
| t | 9.10 | 1 | 10 | 10 |
| a | 8.12 | 00 | 110 | 11 |
| 0 | 7.68 | 01 | 1110 | 010 |
| i | 7.31 | 10 | 1111 | 011 |



## Prefix Codes

- Such codes were invented by Huffman.
- as a term paper at MIT during his PhD.
- had the habit of keeping poisonous snakes as pets.
- Prefix codes are easy to decode.
- No ambiguous decoding possible.
- Faster transmission of frequent data.
- In practice, close to 40-50\% improvement
- We will study Huffman's algorithm during Heaps.


## Learning Outcomes

- Apply tree data structure in relevant applications.
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- Analyze complexity of various operations.

