

CS 2700 PROGRAMMING & DATA STRUCTURES.

WEEK 4 : LISTS

GOALS FOR THE WEEK

- 1) ABSTRACT DATA TYPES (ADTs)
- 2) LIST ADT
- 3) IMPLEMENTING A LIST ADT
- 4) APPLICATIONS .

WHAT IS AN ADT / WHY ADTS?

ADT IS A MATHEMATICAL ABSTRACTION OF DATA.

- FOCUS ON WHAT AND NOT ON HOW.

- example: int data type in C / C++
supports +, -, mod, /, *

details of implementation are hidden / abstracted

- DEFINING AND USING ADTS enables modular programming

BASIC DATA TYPES

int, float, char

bool (C++)

USER DEFINED DATA TYPES

```
struct student { //members }
```

```
struct node { ... }
```

```
struct rectangle { }
```

ABSTRACT DATA TYPES.

Student ADT : what interface does it provide?

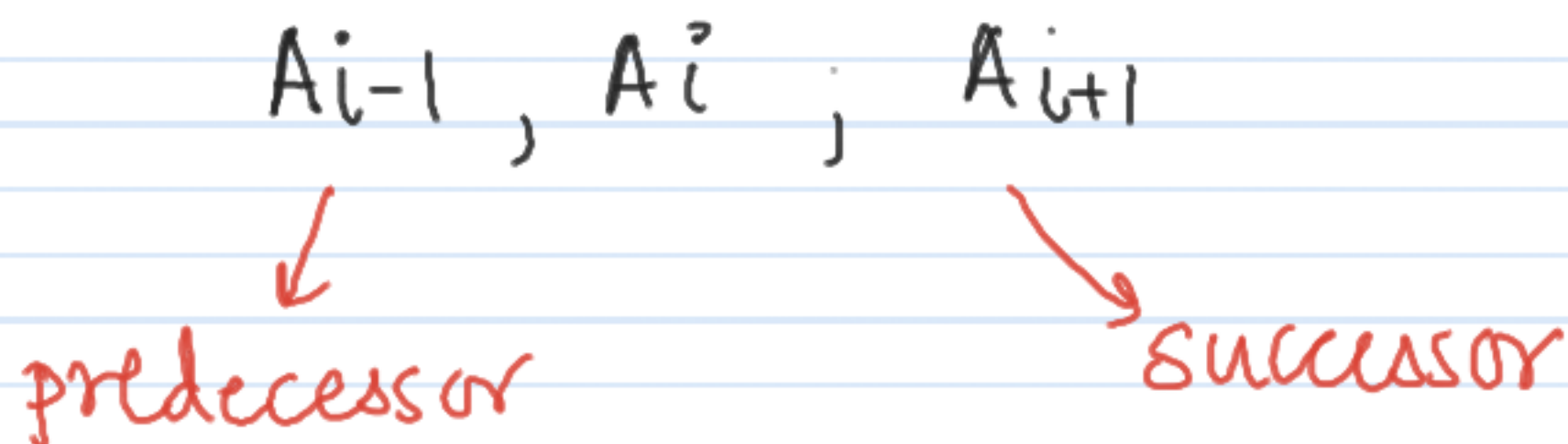
- set Roll No, get Roll No

- set courses, get courses, get Dept, set dept

LIST ADT

$A_1, A_2, A_3, \dots, A_n$

- general list



- operations on list

1. add an element
2. delete an element
3. search for an element
4. print list

LIST ADT.

OPERATIONS :

1. SIZE

2. INSERT ELEMENT val

allow duplicates
insert in the end

3. DELETE val → delete all occurrences

4. FIND val → true if present
false otherwise

5. PRINT

ARRAY BASED IMPLEMENTATION OF LIST

MAINTAIN AN ARRAY INTERNALLY

— HOW BIG?

assume some max size.



free cells in array

INSERT

— maintain index at pos to be inserted
 $O(1)$

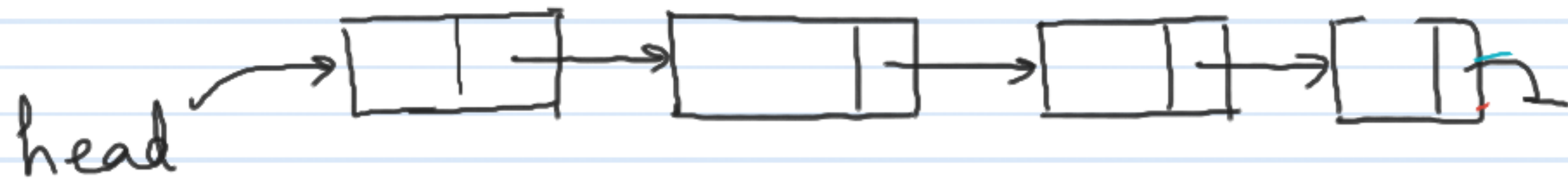
DELETE

— when we see element 'e' move all elements to the right of e by 1 post

improve this to $\Theta(n)$

$\left\{ \begin{array}{l} k \text{ occurrences of element } e \\ \text{each delete takes } O(n), \text{ TOTAL TIME } O(n^2) \end{array} \right.$

LINKED LIST BASED IMPLEMENTATION.



INSERT $O(n)$ since insert happens at end

DELETE $O(n)$ delete needs to go over all
the list

FIND $O(n)$

SIZE $O(n)$ but can be improved to $O(1)$

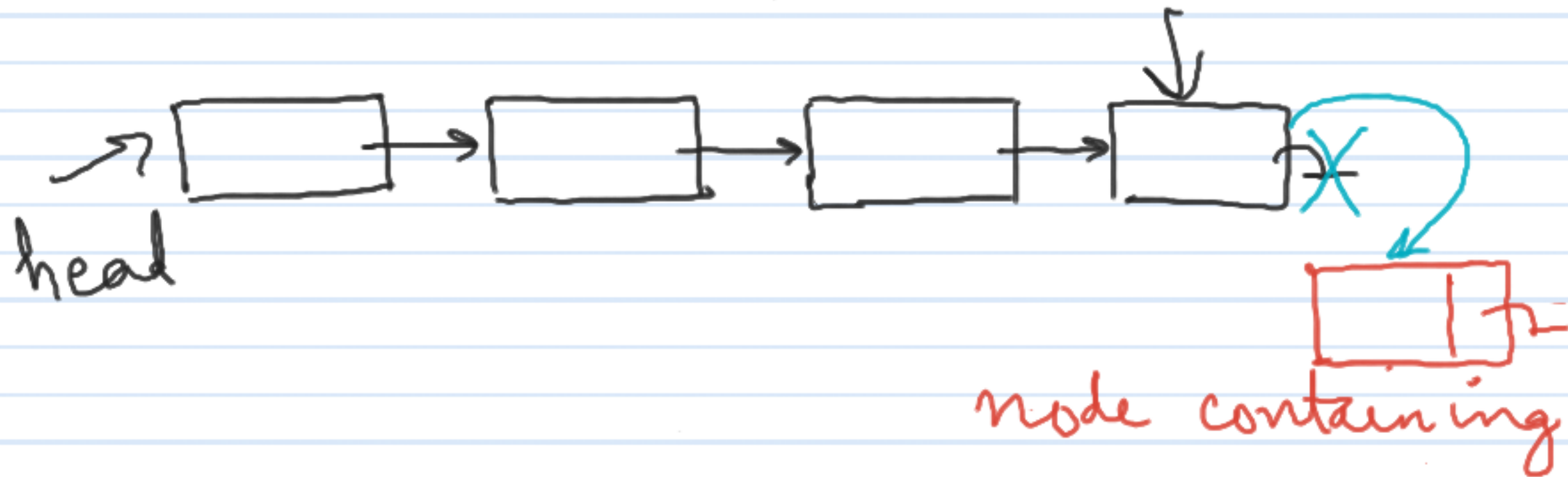
INSERT IN A LINKED LIST

TWO CASES:

(1) INSERT FIRST ELEMENT

head needs to change

(2) INSERT NON FIRST ELEMENT



node containing inserted element

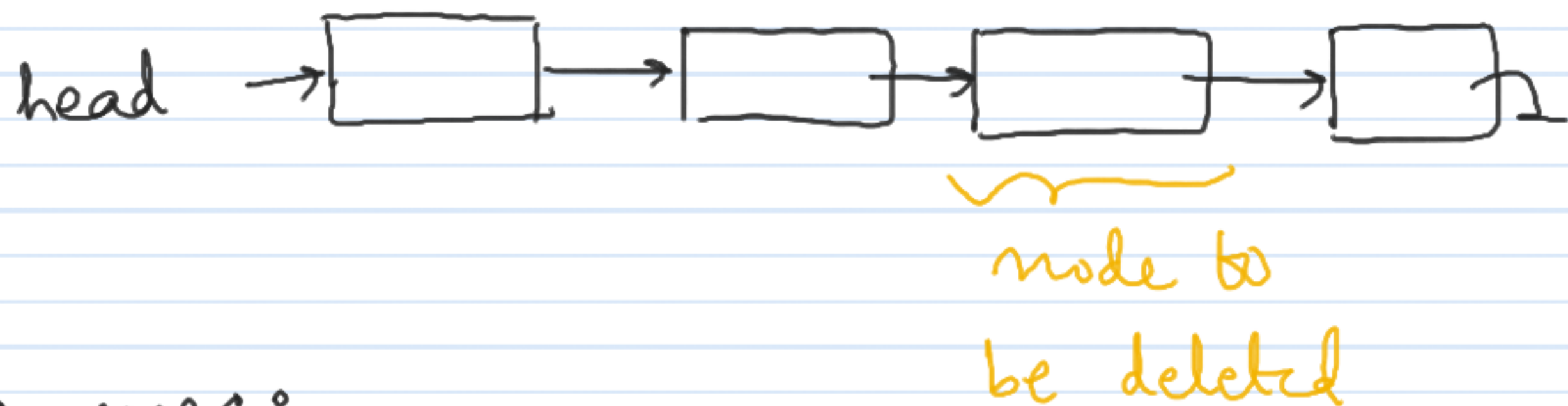
In Both cases:

- we need to get storage on heap (new/malloc)

- set nodes val and next ptr appropriately

These are $O(1)$ operations

DELETE FROM A LINKED LIST.



Two cases:

(1) head needs to be deleted

(2) some internal node needs to be deleted

Maintain a prev pointer, $prev == NULL$ indicates first node needs to be deleted.

Recursive print list

```
void printR ( Node * ptr )
```

```
if ( ptr == NULL )
```

```
    return ;
```

```
else {
```

```
    { ptr->print() ;
```

```
      printR( ptr->next ) ;
```

```
    }
```

```
}
```

simply
prints the
node.

Notes:

(1) If head is "private" then a starter function is needed. (see code)

(2) The running time of the function is $T(n) = T(n-1) + O(1) \approx O(n)$.

} What happens if we switch the order of the 2 statements?