CS6013 Assignment 1

1. Regular Expressions and DFA

Draw DFAs for the following languages (5 + 5 + 5)

- (a) L={ $w \in \{a, b, c, d\}^* | w$ has no repeated letters}
- (b) L={ $w \in \{1,2,3\}^*$ | Number of 2s modulo 2 equal the number of 3s modulo 3}
- (c) L= $\{w \in \{lock, unlock, tryLock\}^* | w$ denotes a sequence of valid lock operations in Java on a lock $\}$.

Bonus: Write the equivalent REs. (10)

2. **CFG**

Write the CFG for the following language: (5 + 5 + 5)

- (a) L={ $w \in \{0,1\}^* | w \text{ contains equal number of 0s and 1s}}.$
- (b) L= $\{w \in \{0,1\}^* | w \text{ contains unequal number of 0s and 1s} \}$.
- (c) L= $\{w \in \{push, pop, top\}^* | w \text{ denotes a sequence of valid stack operations}\}.$

3. Parsing

LL(1) Grammar (30), Parser Implementation (40).

Consider the grammar

stmt ...= id(); | stmt stmt | { stmt } | if (id) stmt else stmt
where stmt is the only non-terminal symbol, stmt is the start symbol, and
id () ; { } if else

is the list of terminal symbols. The terminal symbol id is defined using the regular expression (letter+) where letter is an ascii character in the interval a...z. The grammar generates a subset of the Java statements. Rewrite the grammar into a grammar which is LL(1), and use the rewritten grammar as the basis for implementing a recursive descent parser: write the LL(1) grammar, the FIRST and FOLLOW sets for each nonterminal symbol, and the predictive parsing table, together with an argument that the new grammar is LL(1).